

SOFT-DECISION DECODING OF REED-SOLOMON CODES

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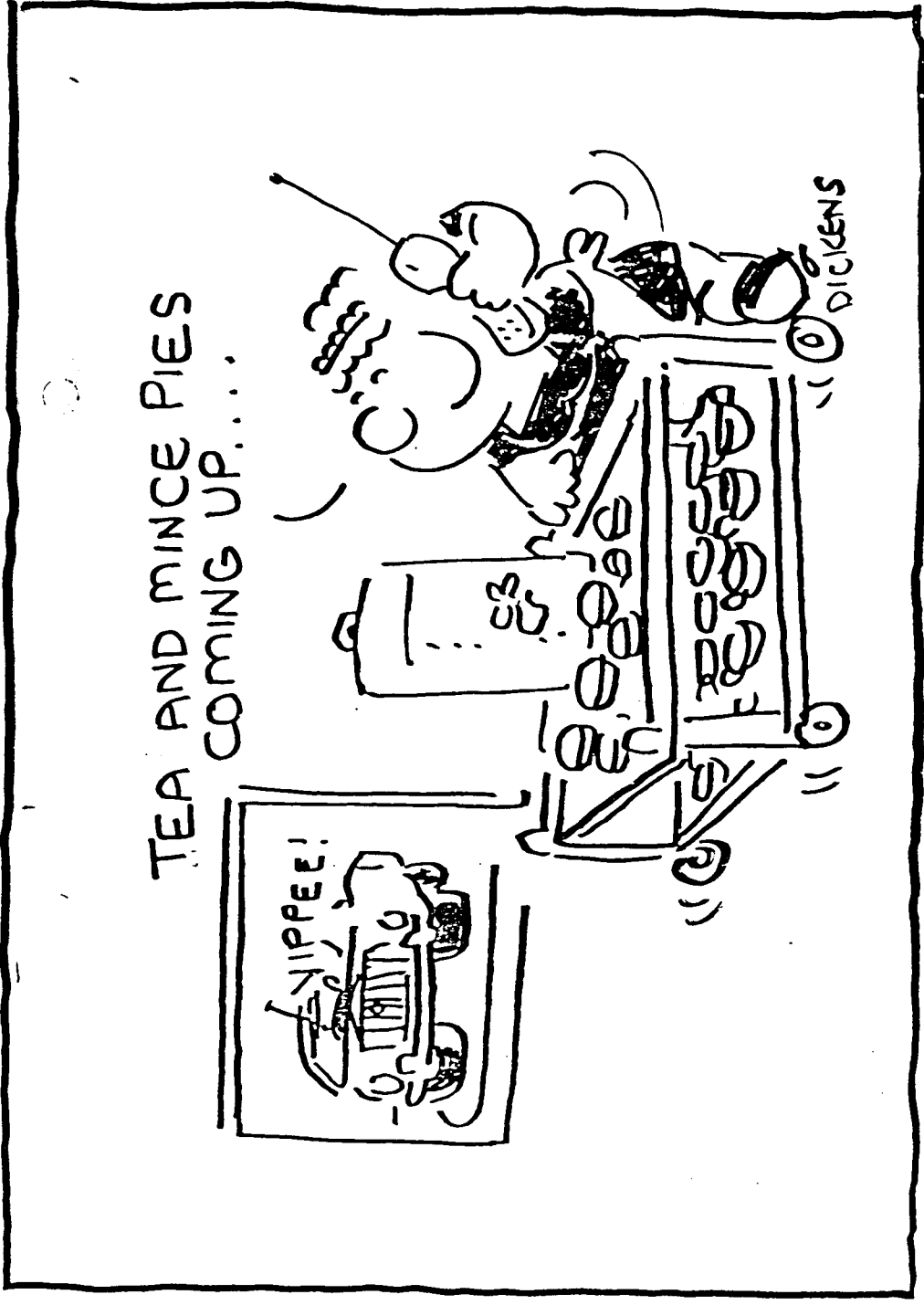
- NON-ALGEBRAIC DECODING
- SOFT-DECISION DECODING
- CODED MODULATION

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SOFT-DECISION DECODING OF RS CODES

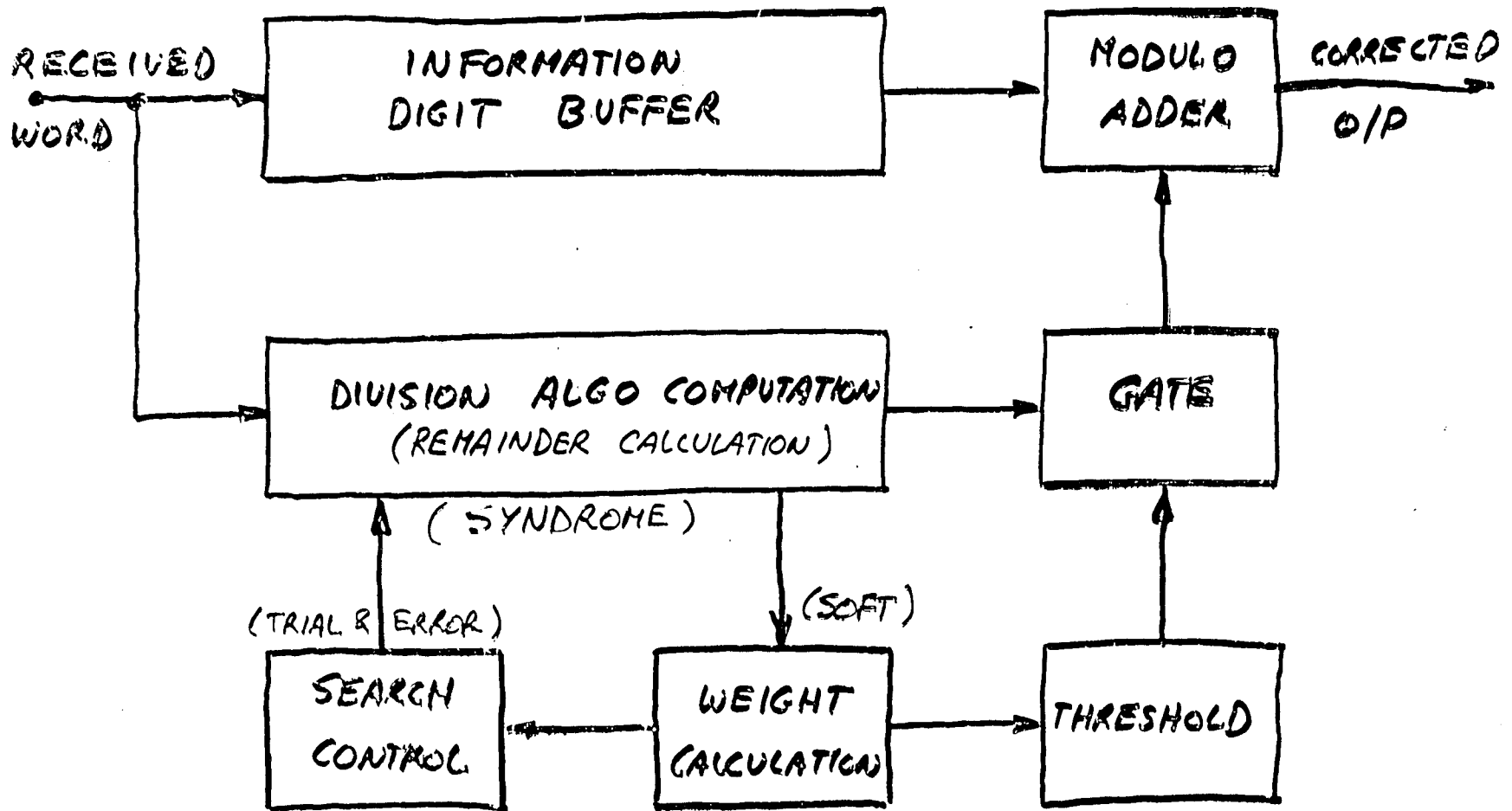
The most effective decoding algorithms for cyclic codes are those, which exploit the code structure and symmetry to the fullest possible extent, are extendable to soft-decision decoding and can offer trade-off between decoding performance and complexity. Minimum weight decoding seems to fulfill these requirements. It is based on the error trapping technique, which is the simplest of the syndrome decoding methods, combined with trial-and-error (systematic search) and step-by-step decoding methods, resulting a very effective reduced-search bounded minimum-distance decoding algorithm. It makes full use of the nature of a cyclic code and since it is essentially a form of minimum-distance decoding, its adaption for soft-decision decoding is quite simple, only requiring changing the metrics to soft (Euclidean) distance from the hard (Hamming) distance. Its complexity is proportional to the square of the block length or less making it very attractive for decoding short and medium length codes. Minimum weight decoding, a burst-and-random-error-correcting algorithm, is also an interpretative and unifying concept which explains the fundamental nature of syndrome decoding methods, and relates apparently distinct algorithms within a coherent framework. It manifests a practical example of an implementation of step-by-step decoding. In the trial-and-error stage of the decoding algorithm there exists an upper bound on the maximum number of trial-and-error positions necessary to be tested (i.e., not all trial-and-error positions need to be tried) which reduces the decoding delay (or latency) and overall complexity quite considerably.

A study of the extension of hard- and soft-decision minimum weight decoding to multi-level cyclic Reed-Solomon codes reveals that if the trial-and-error stage of the basic decoding algorithm is modified, then for example, with the RS (15,9) code, up to 1dB in the hard-decision case and 2dB in the soft-decision case is gained, when bit errors, rather than symbol errors, are corrected. Here the decoding algorithm can offer trade-off between complexity and performance. For example, in the soft-decision case the trial-and-error stage can be much simplified, and therefore complexity reduced, at the price of 1dB loss in the coding gain. The minimum weight decoding algorithm also can be extended to multi-level signalling schemes, such as ASK, m-ary PSK and QAM. On the two-dimensional 16-QAM AWGN channel, for example, simulation results show that the soft-decision minimum weight decoding of RS (15,9) code offers 1.5dB extra coding gain over the hard-decision case and an overall coding gain of 2.5dB. This combined coding and modulation scheme has very low complexity.



MOBILE RADIO
CYCLIC AND CONV. CODES
RS CODES

- MINIMUM WEIGHT DECODING
(CONTINUED DIVISION
ALGORITHM \equiv ERROR TRAPPING)
- HARD-DECISION RS MWD
SYMBOL DECODING
BIT DECODING
- SOFT-DECISION RS MWD
BINARY SIGNALLING
ASK "
QAM "
- SOFT-DECISION RS MLSD/MWD
QAM SIGNALLING



MINIMUM WEIGHT DECODER

THERE IS A YOUNG FELLOW OF BONAS ,
WHO WANTS TO USE CYCLIC DECODERS ;
WHEN HE TRIES SOFT-DECISION ,
FOR BETTER PRECISION ;
THE GAIN TO HIS SYSTEM IS A BONUS



HARD SYMBOL MWD

$$RS(7, 3, 5) \quad t = 2 \quad m = 3$$

$$G(x) \equiv 1 \alpha^4 \alpha^2 \alpha^4 1$$

1	1	1	1	1	1	1	X(x)
1	0	0	1	0	0	0	E(x)
0	1	1	0	1	1	1	Y(x)
	1	α^4	α^2	α^4	1		G(x)
0	0	α^5	α^2	α^5	0	1	$w = 4$
		α^5	α^2	1	α^2	α^5	
0	0	0	0	α^4	α^2	α^4	3
α	α^4			α^4	α	α^6	
α	α^4	0	0	0	α^4	α^3	4
α^6	α	α^4			α^4	α	
α^5	α^2	α^4	0	0	0	1	4
α^4	α^2	α^4	1			1	
1	0	0	1	0	0	0	$w = 2$: ERRORS

0	000
1	001
α	010
α^2	100
$\alpha^3 = \alpha + 1$	011
$\alpha^4 = \alpha^2 + \alpha$	110
$\alpha^5 = \alpha^3 + \alpha^2 = \alpha^2 + \alpha + 1$	111
$\alpha^6 = \alpha^3 + \alpha^2 + \alpha = \alpha^2 + 1$	101

HARD BIT MWD

001 001 001 001 001 001 001	$X(z)$
001 000 000 001 000 000 000	$E(z)$
<hr/>	
000 001 001 000 001 001 001	$Y(z)$
001 110 100 110 001	$G(z)$
<hr/>	
000 000 111 100 111 000 001	$w_{bit} = 8$

111 100 110 000 000 000 001	$w_{bit} = 7$
110 100 110 001 001	
<hr/>	
001 000 000 001 000 000 000	$w_{bit} = 2$, ERRORS

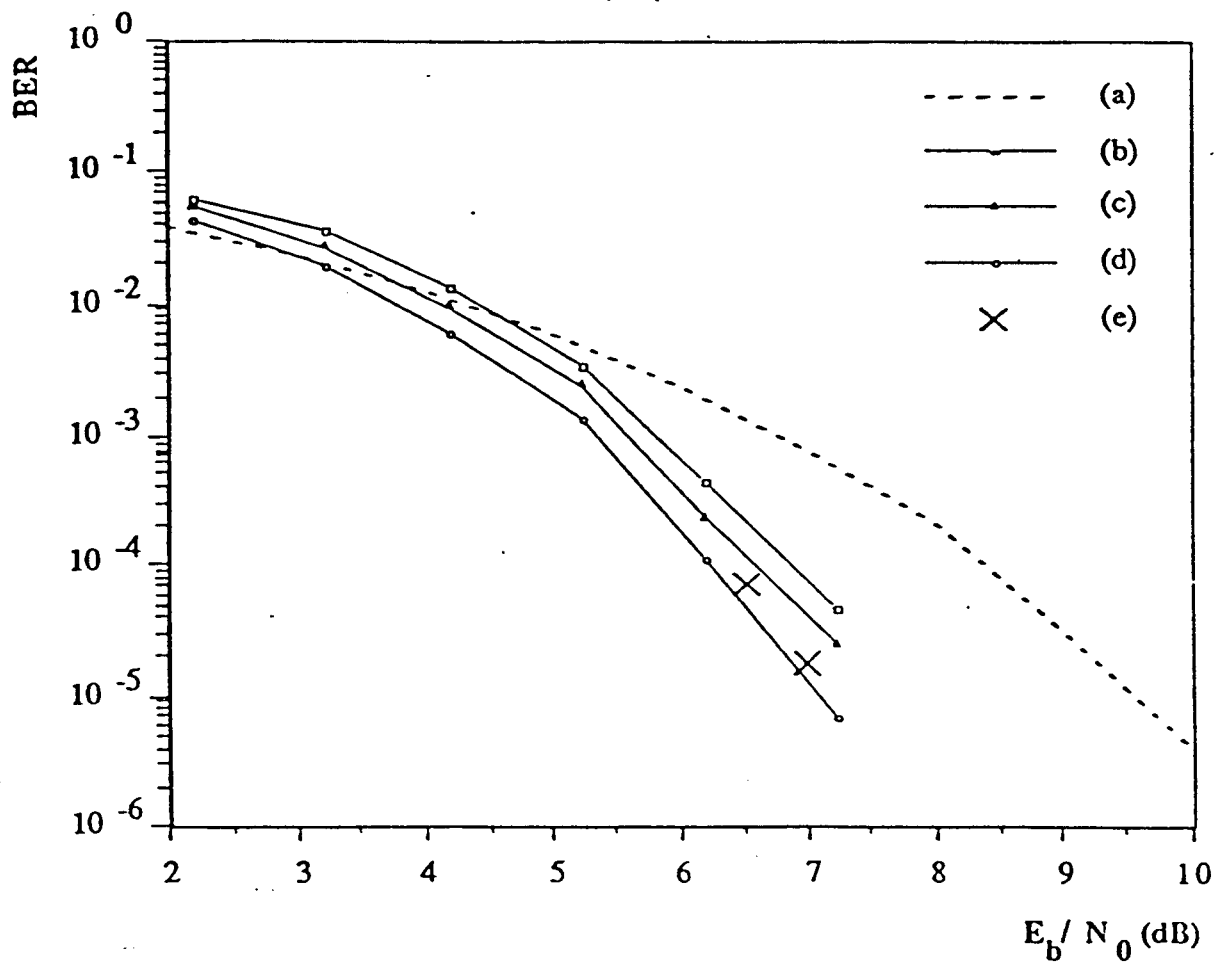
FOR 2 SYMBOL ERRORS :-

$$2 \leq w_{bit} \leq 6$$

WITH THE RS(15, 9, 7) $m=4$, $t=3$ CODE

$$3 \leq w_{bit} \leq 12$$

AND OPTIMUM THRESHOLD IS FOUND TO BE 4



Hard-decision MWD Performance of the RS (15, 9) Code.

(a) uncoded, (b) $t = 3$ (symbol decoding), (c) $t_{bit} = 4$ (bit decoding),

(d) $t_{bit} = 4$ (bit decoding with modified trial-and-error decoding),

(e) hard-decision maximum-likelihood (symbol) decoding (calculated).

TRIAL & ERROR :

IF INITIAL TRAPPING FAILS , CONVERT EACH RECEIVED SYMBOL TO ALL OTHER SYMBOLS , IN TURN , ATTEMPTING TO TRAP AFTER EACH CONVERSION

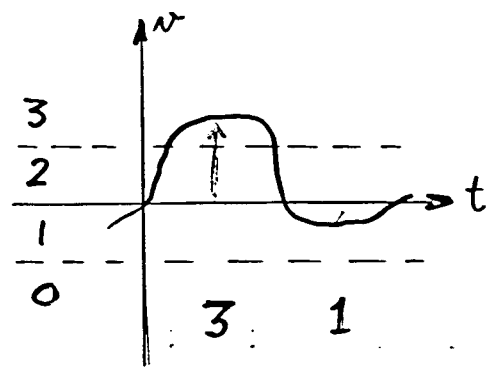
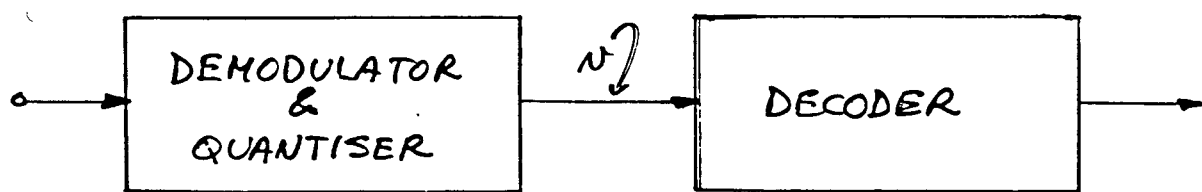
MODIFIED (EXTENDED) TRIAL & ERROR:

IF ABOVE FAILS , CONVERT ALL PAIRS OF RECEIVED SYMBOLS TO SYMBOLS ONE BIT AWAY , IN TURN , AND ATTEMPT TO TRAP

MODIFIED AND SIMPLIFIED TRIAL & ERROR:

AS FOR MODIFIED CASE , BUT INITIAL TRIAL & ERROR CONVERSION IS ONLY TO SYMBOLS ONE AND TWO BITS AWAY

SOFT DECISION DECODING



Q = NUMBER OF QUANTISATION REGIONS (= 4)

SOFT DISTANCE :

BETWEEN DIGITS = $|n_i - n_j|$

BETWEEN WORDS = $\sum_{l=1}^n |n_{il} - n_{jl}|$

EXAMPLE

3102
1230

2132
=> 8

SOFT BIT MWD (BINARY SIGNALS)

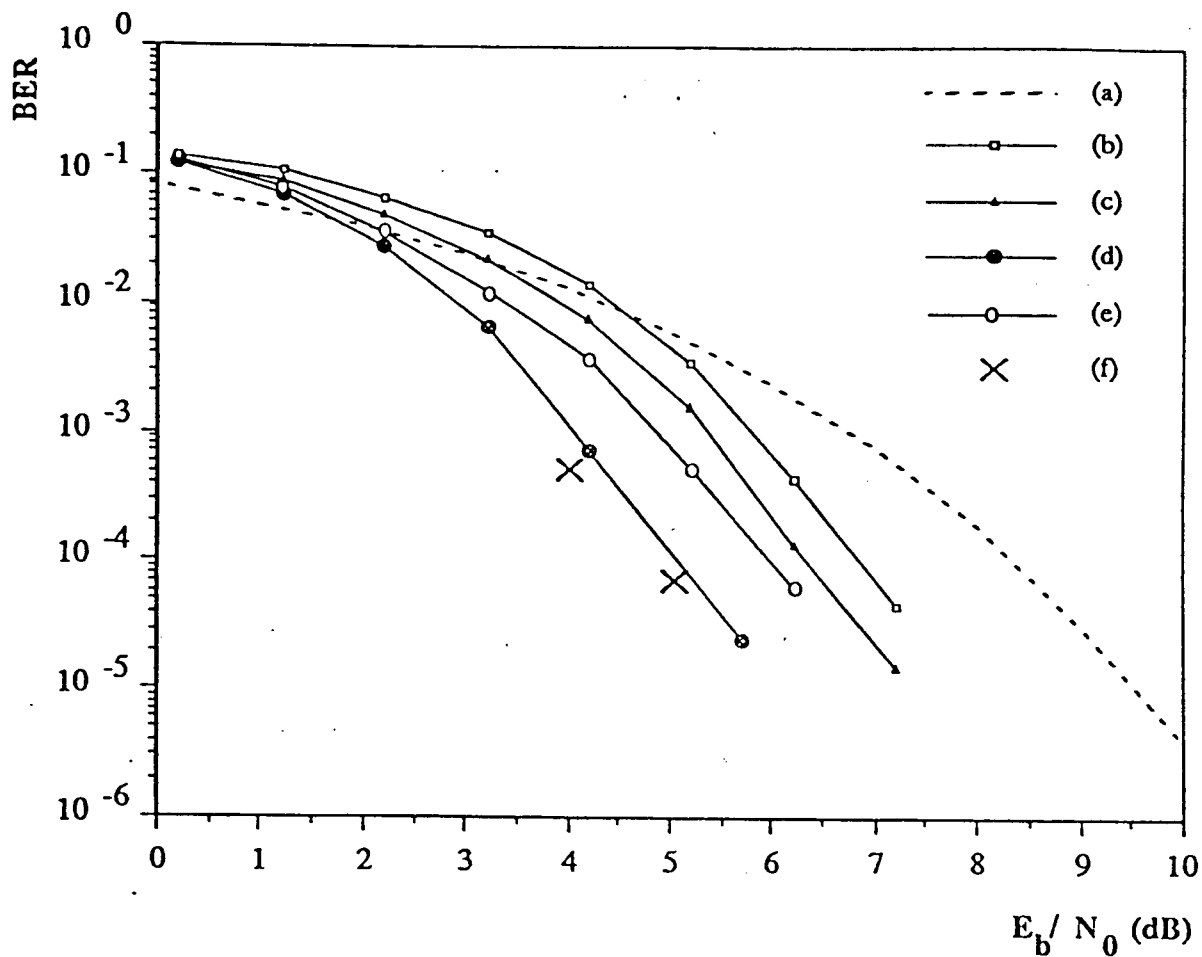
$$t_s = \left\lfloor \frac{d(Q-1)-1}{2} \right\rfloor = 17 \text{ FOR } Q=8$$

007 007 007 007 007 007 007	X(x)	
007 000 000 005 020 000 000	E(x)	
<hr/>		
000 007 007 002 027 007 007	Y(x)	
007 770 700 770 007	G(x)	
<hr/>		
000 000 777 702 757 000 007	$w_{sbit} = 56$	
777 700 007 700 777	$\alpha^5 G(x)$	
<hr/>		
000 000 000 002 750 700 770	33	
070 770 770 070 707	$\alpha^4 G(x)$	
<hr/>		
070 770 000 002 020 770 077	53	
707 070 770 770 070	$\alpha^4 G(x)$	
<hr/>		
777 700 770 002 020 000 007	53	
770 700 770 007 007	G(x)	
<hr/>		
007 000 000 005 020 000 000	14	
	ERRORS	

FOR RS(15, 9, 7) m=4 AND Q=4

$$12 \leq w_{sbit} \leq 84$$

OPTIMUM THRESHOLD IS 42



Soft-decision MWD Performance of the RS (15,9) Code.

(a) uncoded, (b) $t = 3$ (hard symbol decoding),

(c) $t_{s-bit} = 42$ (soft bit decoding),

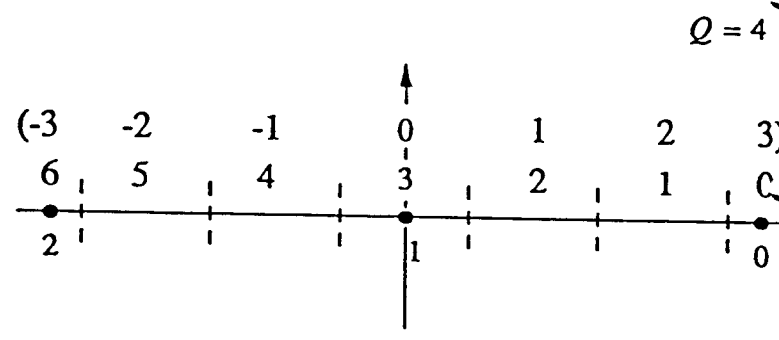
(d) $t_{s-bit} = 42$ (modified trial-and-error decoding),

(e) $t_{s-bit} = 42$ (modified & simplified trial-and-error decoding),

(f) soft-decision maximum-likelihood (symbol) decoding (calculated).

SOFT-DECISION WITH ASK SIGNALLING

- high confidence 2 → 6
- low confidence 2 → 5
- low confidence 1 → 4
- high confidence 1 → 3
- low confidence 1 → 2
- low confidence 0 → 1
- high confidence 0 → 0



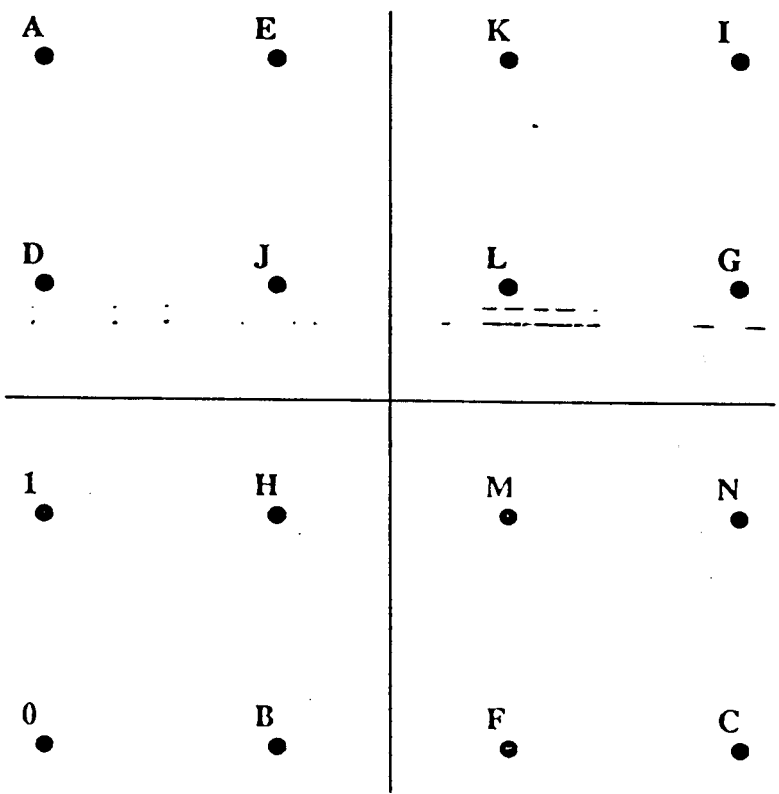
RS (3, 2, 2) IN GF(3)

$$G(z) = z + 2 \equiv 12$$

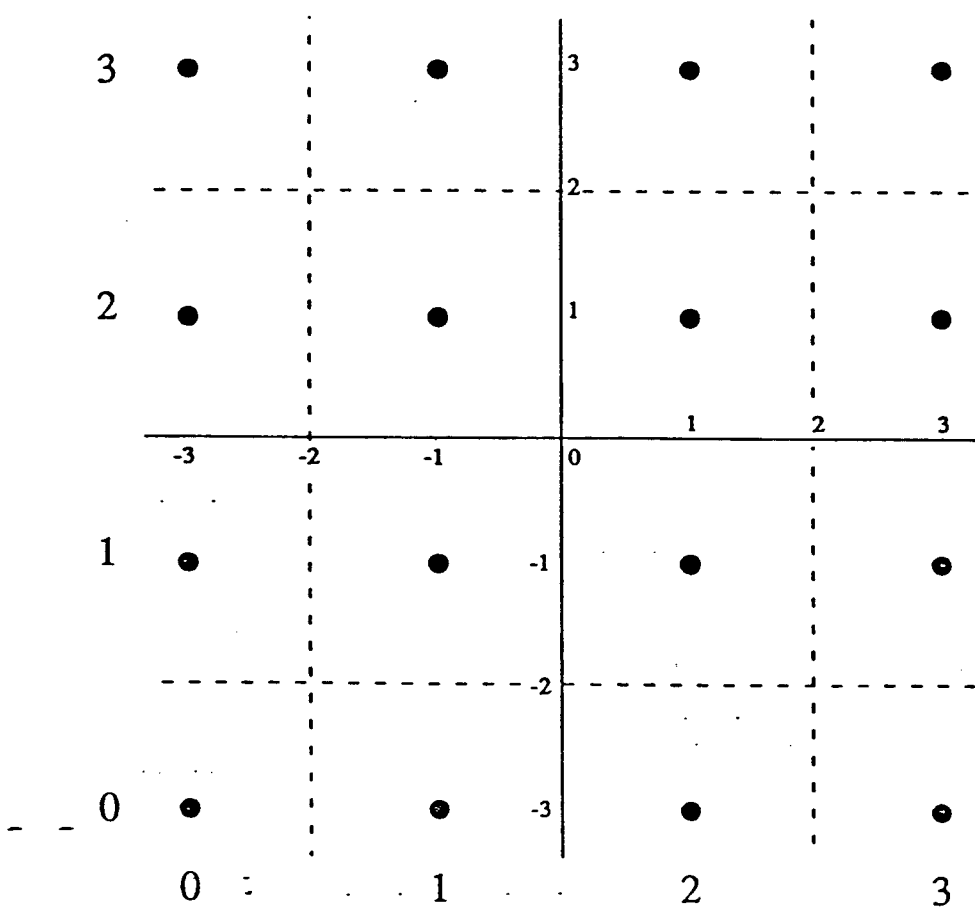
$$t_s = \left\lfloor \frac{d(q-1)-1}{2} \right\rfloor = 2$$

	SOFT WT.	
6 6 6	↓	X
0 2 0		E
<hr/> 6 4 6	16	Y
6 3		2.G
<hr/> 0 1 6	7	(S)
3 6		2.G
<hr/> 3 1 0	4	
3 6		G
<hr/> 0 5 0	5	
6 3		2.G
<hr/> 0 1 3	4	
6 3		2.G
<hr/> 6 1 0	7	
6 3		2.G
<hr/> 0 2 0	2	SOFT ERRORS

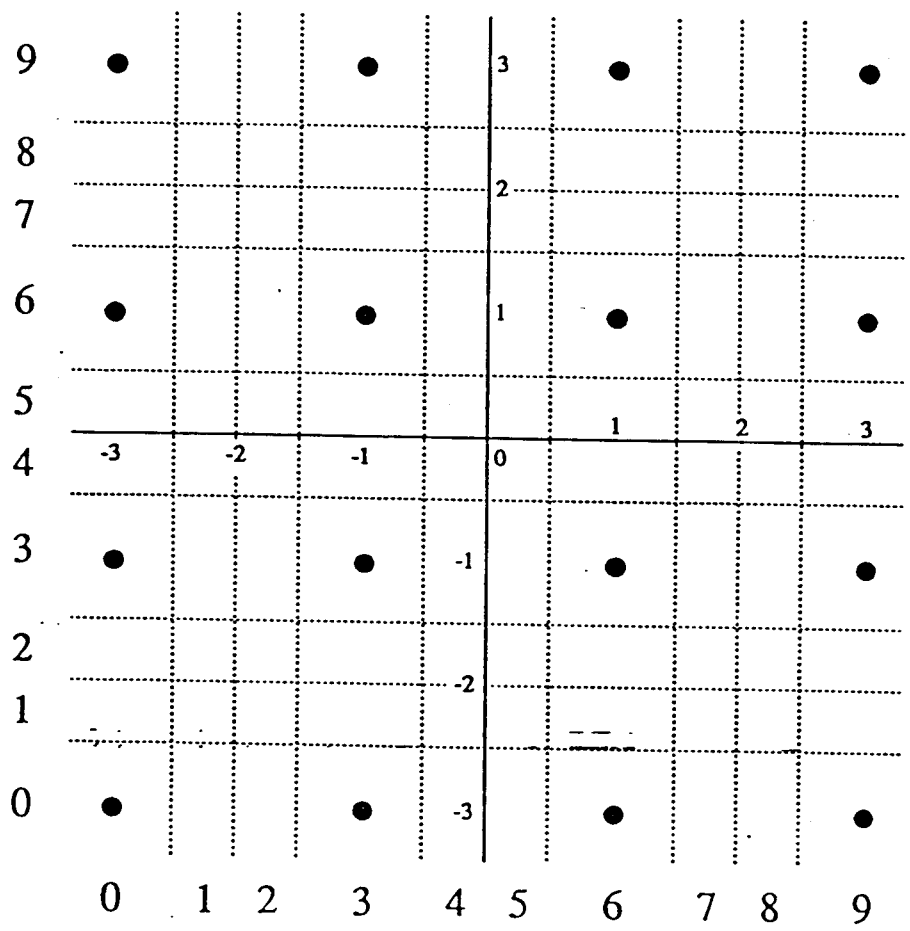
symbol	as a 4-tuple
0	0000
1	0001
A	0010
B	0100
C	1000
D	0011
E	0110
F	1100
G	1011
H	0101
I	1010
J	0111
K	1110
L	1111
M	1101
N	1001



Two-dimensional Gray mapped 16-QAM signal constellation.



Hard-decision quantisation thresholds and corresponding confidence values on the 16-QAM channel.



Soft-decision quantisation thresholds and corresponding confidence values on the 16-QAM channel.

$$Q = 4$$

The MWD algorithm in the two-dimensional QAM channel, for example, would consist of the following steps :

Let the transmitter codeword be c , and the received signal be r at the decoder;

Divide r by the first finite field subtractor (*i.e.*, the appropriate shift of the generator sequence, $G(x)$, multiplied by the appropriate field element; using the corresponding decision symbol in the soft-decision case) thus forming a syndrome (the first finite field subtractor is also the first candidate codeword);

(*) Calculate the Euclidean distances between the symbols of r and the symbols of the candidate codeword, where the Euclidean distance is the root of the sum of the squares of the metrics in each dimension;

Calculate the "weight", w_e , by adding up the Euclidean distances found in (*);

If $w_e \leq t_e$ (threshold determined experimentally for HD and SD)

the candidate codeword is taken as the decoded codeword;

stop;

else Shift $G(x)$ cyclically, and find the next corresponding finite field subtractor with respect to the previous syndrome;

Divide r by the sum of (all) subtractors found so far, *i.e.*, by the new candidate codeword, thus forming a new syndrome;

Go to (*) and continue;

If after a total of $k + n$ shifts of $G(x)$, w_e has never fallen to t_e or less, convert each symbol of r in turn (in the soft-decision case, least confidence symbols first) to all other possible symbols in $GF(q)$, and do the following in each case :

Divide r' (r with a converted symbol) by the first finite field subtractor to form a syndrome. The first finite field subtractor is also the first candidate codeword.

(*) Calculate the Euclidean distances between the symbols of r and the symbols of the candidate codeword;

Calculate the "weight", w_e , by adding up the Euclidean distances found in (*);

If $w_e \leq t_e$ (threshold determined experimentally for HD and SD)

the candidate codeword is taken as the decoded codeword;

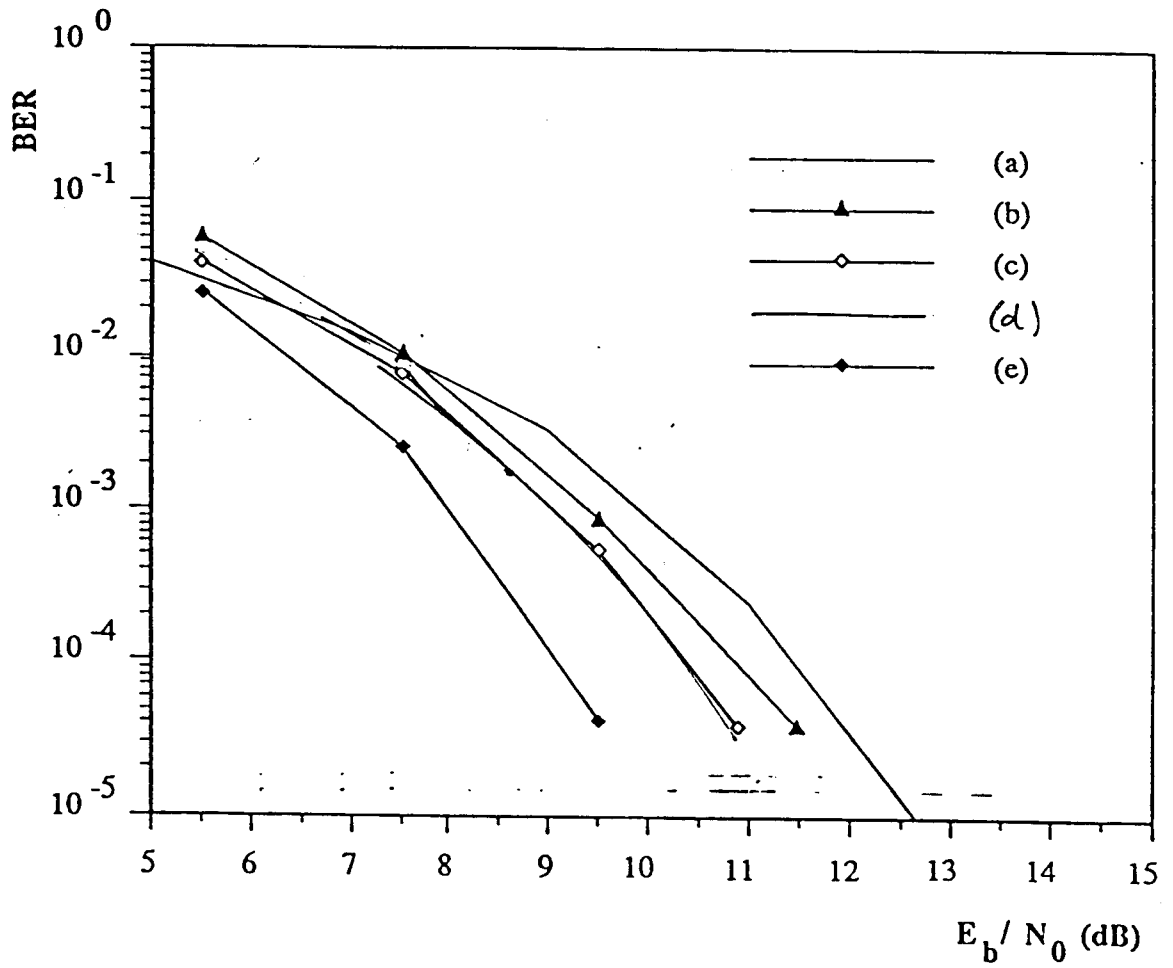
stop;

else Shift $G(x)$ cyclically (a total of $k + n$ times), and find the next corresponding finite field subtractor and syndrome after each shift;

Divide r' by the sum of (all) subtractors found so far, *i.e.*, by the new candidate codeword, thus forming a new syndrome;

Go to (*) and continue.

The algorithm stops when $w_e \leq t_e$ is found, or after all n received symbols have been converted and all the steps for each r' have been performed, in which case no correction is made.

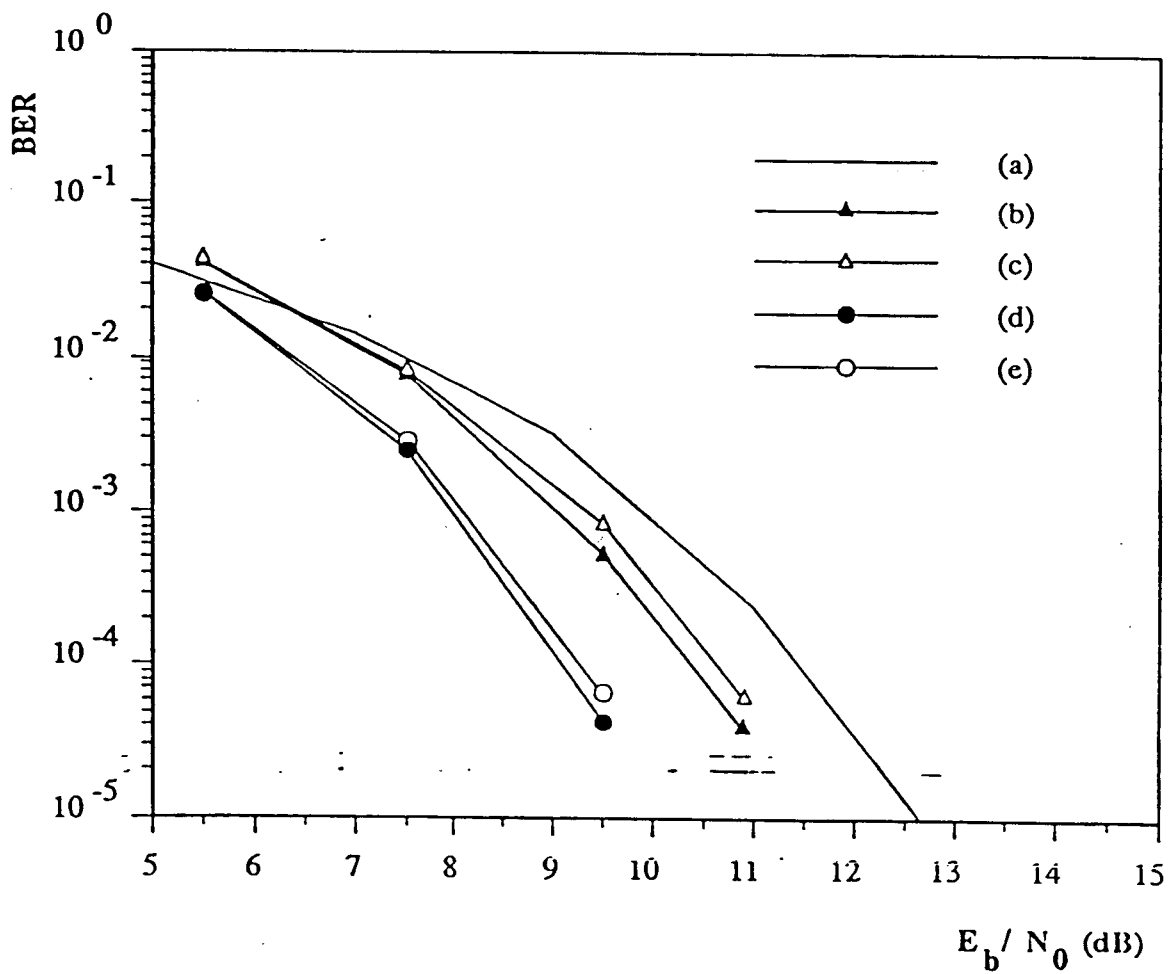


MWD performance of RS (15,9) code on the 16-QAM channel.

(a) uncoded, (b) $t = 3$ (hard symbol decoding),

(c) $t_{e-hard}^* = 4$, (d) RS(15,11) $t_{e-soft}^* = 14$, (e) $t_{e-soft}^* = 16$.

* HARD EUCLIDEAN DISTANCE 1 OR $\sqrt{2}$
IN TRIAL & ERROR STAGE



MWD performance of RS (15,9) code on the 16-QAM channel

- (a) uncoded, (b) $t_{e-hard} = 4$, (c) $t_{e-hard} = 4$ (simplified trial-and-error decoding),
- (d) $t_{e-soft} \cong 16$, (e) $t_{e-soft} \cong 16$ (simplified trial-and-error decoding).

HD EUCLIDEAN DISTANCE 1
IN TRIAL & ERROR

MAXIMUM LIKELIHOOD SEQUENCE DECODING (MLSD) ALGORITHM (TAIT, DORSCH)

BLOCK CODE C
 TRANSMITTED WORD c
 RECEIVED VECTOR r

GENERATE SEQUENCE OF VECTORS x_0, x_1, \dots
 WHERE x_i IS i^{th} MOST LIKELY ESTIMATE OF
 c ; $i \geq 0$:

$$p(r/x_0) \geq p(r/x_1) \geq \dots$$

DECODING IS COMPLETE WHEN $x_i \in C$

COMPLEXITY DEPENDENT ON COVERING
 RADIUS OF CODE (RADIUS OF SPHERE GUARANTEED
 TO ENCLOSE A CODE WORD, CENTRED AT r)

$$(q-1) \sum_{i=0}^S \binom{M}{i} \text{ ITERATIONS}$$

GUIDED BLIND WALK AMONG
 SPARSE CODE WORDS!

COMBINED MLSD AND BOUNDED DISTANCE
DECODING (eg MWD).

GUIDED SIGHTED WALK !

MWD WITH $0 \leq \epsilon \leq t$ ERROR CORRECTION

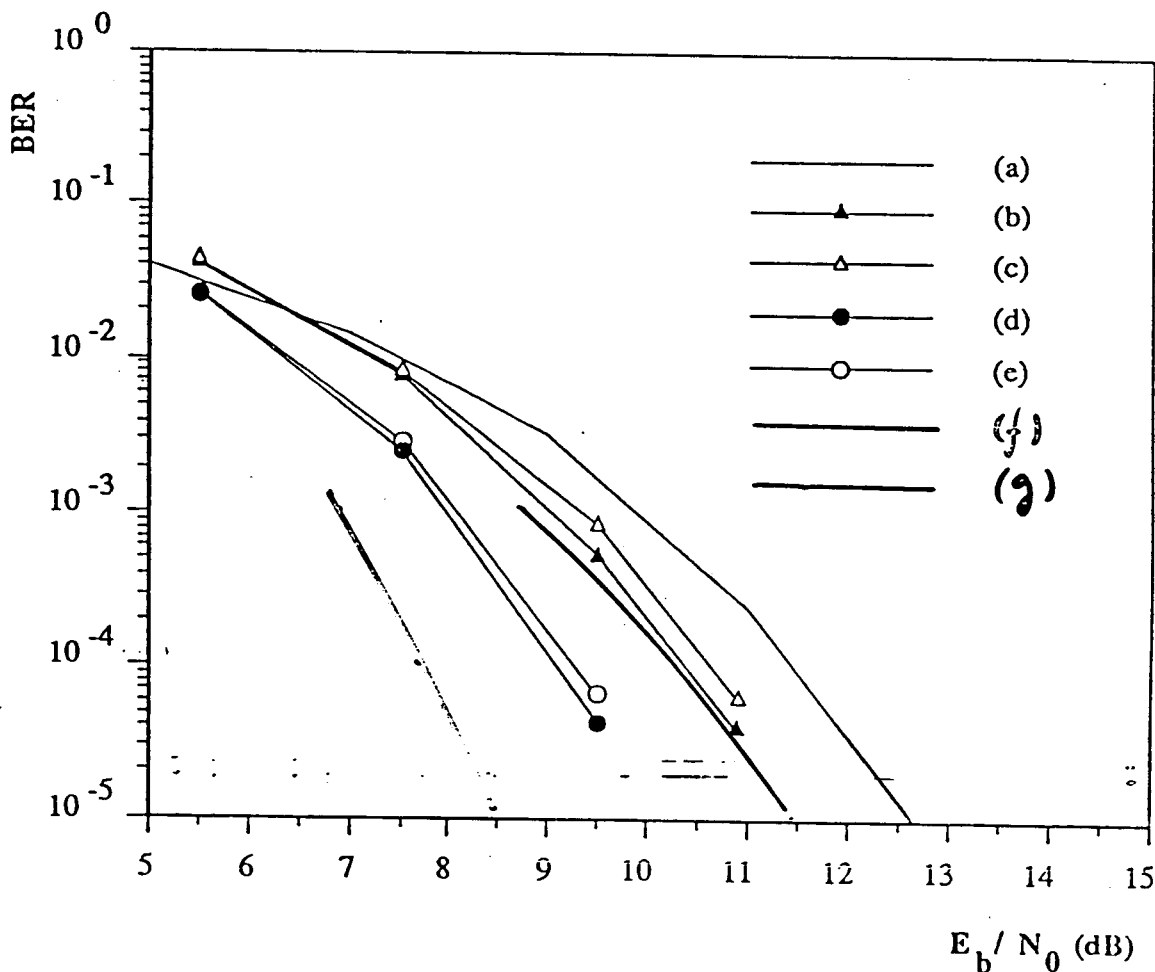
MLD IF $d \{ x_i, r \} \leq t$

OR O/P BEST x_i AFTER FIXED
NUMBER (LARGE) OF ITERATIONS

\equiv MLSD IF $\epsilon = 0$

COMPLEXITY $\propto 1/\epsilon$

CAN BE ADAPTED FOR ANY SD METRIC,
SUCH AS EUCLIDEAN DISTANCE IN A MULTI-
LEVEL SIGNAL CONSTELLATION

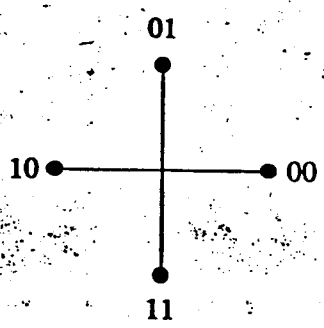


MWD performance of RS (15,9) code on the 16-QAM channel.

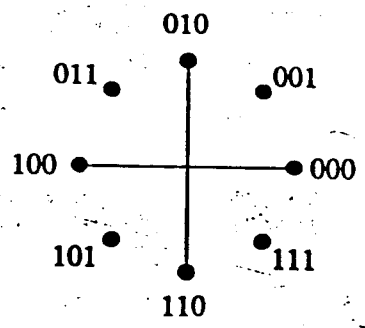
(a) uncoded, (b) $t_{e-hard} = 4$, (c) $t_{e-hard} = 4$ (simplified trial-and-error decoding),

(d) $t_{e-soft} \cong 16$, (e) $t_{e-soft} \cong 16$ (simplified trial-and-error decoding).

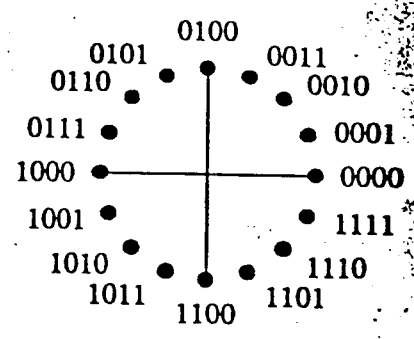
(f) MLSD/MWD (g) 8-QAM UNCODED



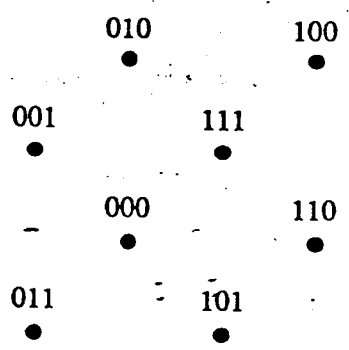
4-PSK



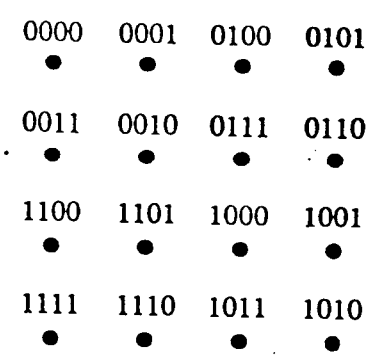
8-PSK



16-PSK



8-QAM



16-QAM

Signal Set Constellations